

DESIGNING FOR CONCURRENCY USING MESSAGE PASSING

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Designing For Concurrency Using Message Passing

- Why use message passing?
- Message passing frameworks
- Design Guidelines
- Examples



Why use message passing?



The Hook

Eliminate all your concurrency bugs!



You can:



You can:

• Eliminate explicit synchronization



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- Eliminate OS-level deadlock



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- Eliminate explicit synchronization
- Eliminate OS-level deadlock
- Eliminate data races



The Cost

In exchange, you must:





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• Use the framework for all synchronization





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- Avoid shared mutable state





In exchange, you must:

- Use the framework for all synchronization
- Avoid shared mutable state
- Change your design approach





Problems still remain:

Still potential for race conditions
 ⇒ Which message arrives first?



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- Still potential for message deadlock
 - \Rightarrow Two elements expecting messages from each other



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 ⇒ Which message arrives first?
- Still potential for message deadlock
 - \Rightarrow Two elements expecting messages from each other
- The potential for dropped messages is a **new** problem
- Still not a silver bullet



Message Passing Frameworks



What is a Message Passing Framework?

Tooling that manages the delivery of messages between independent elements



Developers can focus on:



Developers can focus on:

• the content,



Developers can focus on:

- the content,
- the **targets**, and



Developers can focus on:

- the **content**,
- the **targets**, and
- the processing



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of messages, rather than the **delivery mechanism**.



Microservices

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The API for each microservice is the messages you can send it, and the messages it sends.



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They may also guarantee that all messages for a given microservice will be handled on the same thread.

They may give you control over which microservices share a thread.



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Ideally the microservice code will be isolated from that choice.



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Abstract the boiler plate in a wrapper library or tool.



Design Guidelines





Each microservice should be focused on one task



Independence

Microservices should not overlap in their functionality



Value Types

Messages should be value types



Avoid blocking

Replace blocking calls with separate microservices



State Machines

Each microservice can be modelled as a single-thread state machine



Examples



Example: Dining Philosophers



Example: Dining Philosophers

- N philosophers sat round a table sharing a large plate of rice.
- Each philosopher wants to alternate between **thinking** and **eating**.
- There are **N chopsticks** around the table, with one between every pair of philosophers.
- All philosophers start thinking.
- After a random length of time thinking, a philosopher will try and eat.
- To eat, a philosopher must pick up the chopsticks either side. A philosopher cannot eat without both chopsticks.
- After a random length of time eating, a philosopher will put down both chopsticks and think.



Actors for Dining Philosophers

We can create multiple actors:

- An actor for each philosopher
- An actor for I/O
- An actor for Timing
- An actor for the "server" responsible for managing the chopsticks



The I/O Actor

The I/O actor just waits for incoming messages with text to print out.

```
using OutputMessage = std::string;
```

```
void I0::OnMessage(OutputMessage msg) {
   std::cout << msg << '\n';
}</pre>
```



The Timing Actor

The **Timing** actor is more involved. Incoming messages specify a time and another actor:

```
struct TimerRequest {
   std::chrono::steady_clock time;
   messaging::sender target;
};
```

The **Timing** actor sends a **TimerExpired** message to the **target** at the appropriate time.



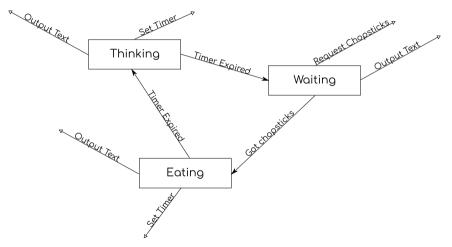
The Philosopher Actors

Each philosopher has 3 states: **thinking**, **waiting for chopsticks** and **eating**.

In each state we can process a particular set of incoming messages.



Philosopher States





A Thinking Philosopher

When the philosopher starts thinking, then we send a message to the **I/O** actor to say so, and to the **Timing** actor specifying when they're done thinking:

```
void Philosopher::OnThinking() {
   auto thinking_time = random_duration();
   io.Send(name + " thinking");
   timing.Send(Clock::now() + thinking_time);
}
```



A Thinking Philosopher

When the timer expires, the philosopher is waiting for chopsticks:

void Philosopher::OnMessage(TimerExpired) {
 if(current_state == states::thinking) {
 io.Send(name + " finished thinking");
 SetState(states::waiting_for_chopsticks);
 }



A Hungry Philosopher

When the philosopher starts waiting for chopsticks, they must be requested from the **server**, after sending a message to the **I/O** actor to be displayed:

void Philosopher::OnWaitingForChopsticks() {
 io.Send(name + " waiting for chopsticks");
 server.Send(RequestChopsticks{
 left_chopstick, right_chopstick,
 GetSelf()});



The philosopher gets chopsticks

When the philosopher gets chopsticks they start eating

void Philosopher::OnMessage(GotChopsticks) {
 io.Send(name + " received chopsticks");
 SetState(states::eating);
}



An Eating Philosopher

The philosopher eats for a random length of time, just like **thinking**, so we send a message to the **Timing** actor specifying when they're done eating:

```
void Philosopher::OnEating() {
   auto eating_time = random_duration();
   io.Send(name + " eating");
   timing.Send(Clock::now() + eating_time);
}
```



A Full Philosopher

When the philosopher is done eating they hand in the chopsticks and resume thinking.

void Philosopher::OnMessage(TimerExpired) {
 if(current_state == states::thinking) {
 // as before
 } else if (current_state == states::eating) {
 io.Send(name + " finished eating");
 server.Send(ReturnChopsticks{
 left_chopstick, right_chopstick});
 SetState(states::thinking);
 }
}



The Server Actor: Chopstick Requests

If chopsticks are requested, the server must hand them out, or add the philosopher to a waiting list:

void Server::OnMessage(RequestChopsticks request){
 if(ChopsticksAvailable(request)){
 UseChopsticksAndNotify(request);
 } else {
 waiting_list.push_back(request);
 }



The Server Actor: Returned Chopsticks

If chopsticks are returned, mark them free and check for waiting philosophers:

void Server::OnMessage(ReturnChopsticks request){
 MarkChopsticksAvailable(request);
 for(auto req=waiting_list.begin();
 req != waiting_list.end();) {
 if(ChopsticksAvailable(*req)) {
 UseChopsticksAndNotify(*req);
 it = waiting_list.erase(it);
 } else ++it;
 }
}



The Server Actor: Using Chopsticks

Notifying philosophers about chopsticks is just another message:

void Server::UseChopsticksAndNotify(
 RequestChopsticks request) {
 chopstick_free[request.left_chopstick] = false;
 chopstick_free[request.right_chopstick] = false;
 request.philosopher.Send(GotChopsticks{});



Example: A robot control system



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Moving a robot requires coordinating movements of multiple parts, each with their own actuator.

We want the control code for each actuator to run concurrently with the coordination code and the code for other actuators.



Example: Moving a robot

We have actors for:

- The central coordinator
- Each actuator
- Timing



The central coordinator

For each large-scale position the robot needs to be in, send messages to the actuators with the steps:

void Coordinator::OnMessage(SetPositionFoo msg) {
 SetState(states::waiting_for_position_foo);
 actuator1.Send(SetPosition{Ac1FooPos, msg.target_time});
 actuator2.Send(SetPosition{Ac2FooPos, msg.target_time});
 actuator3.Send(SetPosition{Ac3FooPos, msg.target_time});
 actuator4.Send(SetPosition{Ac4FooPos, msg.target_time});



Actuator actors

Each actuator has hardware-specific code to control it.

```
void Actuator1::OnMessage(SetPosition msg) {
  hardware.StopMoving();
  auto curpos = hardware.GetCurrentPosition();
  auto delta = msq.position - curpos;
  auto time_delta = msg.target_time - Clock::now();
  if(delta) {
    hardware.StartMoving({.speed = delta/time_delta});
    timing.Send(TimerRequest{msg.target_time, GetSelf()});
```



Actuator actors

Notify the controller when target reached:

void Actuator1::OnMessage(TimerRequest) {
 hardware.StopMoving();
 controller.Send(Actuator1ReachedPosition{});
}



The central coordinator

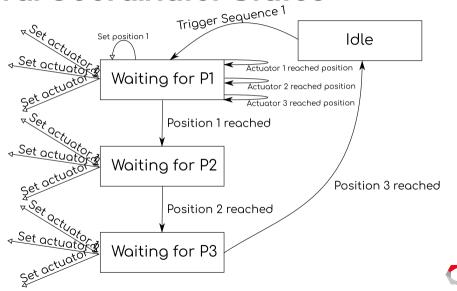
For movement through a sequence of positions, add intermediate states:

- Waiting for Position 1
- Waiting for Position 2
- Waiting for Position 3

Add messages for setting those positions, and notifying when they are reached.



Central Coordinator States



woven

Starting a movement sequence

Start a movement sequence by moving to the first position:

void Coordinator::OnMessage(TriggerMotionSequence1 msg) {
 auto position1_time = Clock::now() + position1_duration;
 SetState(states::waiting_for_position1);
 GetSelf().Send(SetPosition1{position1_time});



Waiting to reach a position

As each actuator reaches position, check overall robot state:

```
void Coordinator::OnMessage(Actuator1ReachedPosition) {
  switch(state) {
  case states::waiting_for_position1:
    if(ActuatorsReachedPosition1()) {
      GetSelf().Send(Position1Reached{}):
   break:
  // code for other states
```



Moving to the next position

When the robot is in the right position, move to the second position:

void Coordinator::OnMessage(Position1Reached) {
 auto position2_time = Clock::now() + position2_duration;
 SetState(states::waiting_for_position2);
 GetSelf().Send(SetPosition2{position2_time});



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Making the state machine table driven can make this easier.

This is just classic state machine processing.





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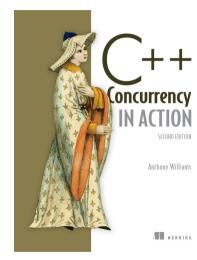


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- Messages should be value types
- Replace blocking calls with separate microservices
- Each microservice can be modelled as a single-thread state machine

My Book



C++ Concurrency in Action Second Edition

Covers C++17 and the first Concurrency TS C++20 Addendum coming soon!

cplusplusconcurrencyinaction.com



Questions?



Message passing frameworks

Some Choices:

- ZeroMQ
- CAF The C++ Actor Framework
- The actor framework from my book
- You can write your own.

