# The Future of Concurrency in C++

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The Future of Concurrency in C++

- Multithreading Support in C++0x
- Existing proposals for TR2
- Beyond TR2

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# Multithreading Support in C++0x

- The Standard now acknowledges the existence of multi-threaded programs
- New memory model
- Support for thread-local static variables
- Thread Support Library
  - Threads
  - Mutexes
  - Condition Variables
  - One time initialization
  - Asynchronous results futures

# C++0x Thread Library and Boost

- Two-way relationship with Boost
  - Proposals for multithreading heavily influenced by Boost. Thread library
  - Boost 1.35.0 Thread library revised in line with C++0x working draft

- Define the rules for making data visible between threads
- Atomics are generally for experts only
- If you correctly use locks, everything "just works"

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There are two critical relationships between operations:

- Synchronizes-with relation
  - Store-release synchronizes-with a load-acquire
- Happens-before relation
  - A sequenced before B in a single thread
  - A synchronizes-with B
  - A happens-before X, X happens-before B

A data race occurs when:

- Two threads access non-atomic data
- At least one access is a write
- There is no happens-before relation between the accesses

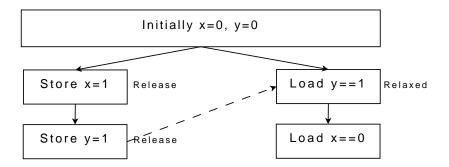
A lot of multithreaded programming is about avoiding data races

# Memory Ordering Constraints

Sequential Consistency

- Single total order for all SC ops on all variables
- default
- Acquire/Release
  - Pairwise ordering rather than total order
  - Independent Reads of Independent Writes don't require synchronization between CPUs
- Relaxed Atomics
  - Read or write data without ordering
  - Still obeys happens-before

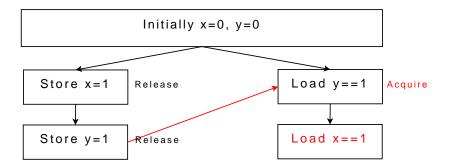
# Relaxed Ordering



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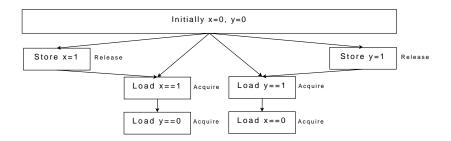
#### Acquire-Release Ordering



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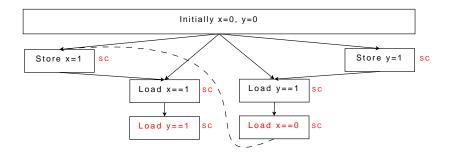
#### Acquire-Release Ordering



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## Sequentially Consistent Ordering



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#### Basic interface for atomics

- atomic\_flag
  - Boolean flag
  - Must be lock-free
- Atomic integeral types e.g. atomic\_char, atomic\_uint, atomic\_llong
  - Includes arithmetic operators such as a++, and a |=5
  - Operators return underlying type by value, not reference
  - May not be lock-free use a.is\_lock\_free() to check
- atomic\_address
  - Represents a void\*
  - May not be lock-free use a.is\_lock\_free() to check
- Free functions for C compatibility

- atomic<T>
  - derived from atomic\_T for built-in integral and pointer types
- works with "trivially default constructible and bitwise equality comparable" types
  - Lock-free where possible

# Compare and Swap

- Generally put in loop
  - Spurious failure
  - Other thread may change value anyway

```
atomic<int> a;
int desired;
int expected=a;
do
{
    desired=function(expected);
}
while(!a.compare_swap(expected,desired));
```



- Per-object fences: a.fence(memory\_order)
   RMW op which writes same value back.
- Global fences with atomic\_global\_fence\_compatibility object (of type atomic\_flag)

std::thread t(func,arg1,arg2);

- std::bind semantics

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std::thread t(func); t.join();

A thread can only be joined once.

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# Detaching a Thread

• Explicitly:

```
std::thread t1(func);
t1.detach();
```

Implicitly:

```
{
    std::thread t2(func);
} // destructor of t2 calls detach()
```

#### Transferring Ownership

- At most one std::thread object per thread.
- Thread objects are movable
  - Can return std::thread from functions

std::thread start\_process(some\_args);

• Can store std::thread objects in standard containers

std::vector<std::thread> vec;

vec.push\_back(std::thread(some\_func));

• Can use t.joinable() to determine if an object has an associated thread.

# Identifying Threads

- Every thread has a unique ID
- Thread IDs represented by instances of std::thread::id
  - Value Type: copyable, usable in comparisons
  - Non-equal values form a total order
  - Can be used as keys in associative containers and unordered associative containers
  - Can be written to an output stream
  - Default constructed ID is "Not any Thread".

# **Obtaining Thread IDs**

- std::this\_thread::get\_id() returns the ID of the current thread
- t.get\_id()

Returns the ID of the thread associated with the std::thread instance t

#### Mutexes

There are four mutex types in the current working paper:

- std::mutex
- std::recursive\_mutex
- std::timed\_mutex
- std::recursive\_timed\_mutex

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# Locking

- lock() and unlock() member functions are public
- Scoped locking:
  - std::lock\_guard template
  - std::unique\_lock template
    - movable, supports deferred locking, timed locking
    - can itself be used as a "mutex".
- Generic lock() function
  - Allows locking of more than one mutex without deadlock

#### Condition Variables

- Two types of condition variables:
  - std::condition\_variable
  - std::condition\_variably\_any
- The difference is the lock parameter to the wait functions:
  - void condition\_variable::wait( unique\_lock<std::mutex>& lock);
  - template<typename lock\_type>
    void condition\_variable\_any::wait(
     lock\_type& lock);

#### Condition Variables and Predicates

- Condition variables are subject to spurious wakes
- Correct usage requires a loop:

```
std::unique_lock<std::mutex> lk(some_mutex);
while(!can_continue())
{
    some_cv.wait(lk);
}
```

• Predicate version makes things simpler:

```
std::unique_lock<std::mutex> lk(some_mutex);
some_cv.wait(lk,&can_continue);
```

#### Timed waits with condition variables

• The overload of condition\_variable::timed\_wait() that takes a duration is particularly error-prone:

```
while(!can_continue())
{
    some_cv.timed_wait(lk,std::milliseconds(3));
}
```

This may actually be equivalent to just using wait(), in the event of spurious wake-ups

• The predicate overload avoids this problem:

One-time Initialization

#### Provided by std::call\_once

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```
std::once_flag flag;
```

```
std::call_once(flag,some_function);
// calls some_function()
```

std::call\_once(flag,some\_other\_function,arg1,arg2);
// calls some\_other\_function(arg1,arg2)

- std::bind semantics again

#### Lazy initialization of class members

```
class X
ł
    some_resource_handle h;
    std::once_flag flag;
    void init_resource();
public:
    X():h(no resource){}
    void do_something()
    ł
        std::call_once(flag,&X::init_resource,this);
        really_do_something(h);
    }
};
```

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#### Thread-local static variables

- Not yet in WP: N2545 by Lawrence Crowl
- Each thread has its own instance of the variable
- Use the thread\_local keyword: static thread\_local int x;
- Any variable with static storage duration can be declared thread\_local:
  - Namespace-scope variables
  - static data members of classes
  - static variables declared at block scope
- thread\_local variables can have constructors and destructors.

# Asynchronous Value Computation

- Not yet in WP: N2561
  - Deltef Vollman, Howard Hinnant and myself
- Value is result of a task running on another thread.
- No control over how or when value is computed by recipient.
- Answer to how to return a value from a thread.

Two templates for futures:

- std::unique\_future<T> like std::unique\_ptr<T>
  - sole owner
  - read once (move)
- std::shared\_future<T> like std::shared\_ptr<T>
  - multiple owners
  - can be read multiple times (copy)
- Can move a std::unique\_future<T> into a
   std::shared\_future<T>

#### Getting the values: std::unique\_future<T>

#### • R move()

- blocks until ready
- throws if already moved
- throws if future has a stored exception
- o bool try\_move(R&)
  - returns false if not ready() or already moved.
- State query functions:

is\_ready(), has\_value(), has\_exception(), was\_moved()

• Wait for ready:

```
wait(), timed_wait()
```

#### Getting the value: std::shared\_future<T>

- R const& get() operator R const&()
  - Blocks until ready
  - Returns reference to stored value
  - Throws if future has a stored exception
- bool try\_get(R&)
- State Query functions: is\_ready(), has\_value(), has\_exception()
  - No was\_moved() has the result can't be moved
- Wait for ready: wait(), timed\_wait()

Two ways of generating asynchronous values:

- std::packaged\_task<T>
  - value is the result of a function call
- std::promise<T>
  - explicit functions for populating the value

# Packaged Tasks

- A std::packaged\_task<T> is like std::function<T()> it wraps any function or callable object, and invokes it when std::packaged\_task<T>::operator() is invoked.
- Return value populates a std::unique\_future<T> rather than being returned to caller
- Simplest way to get the return value from a thread

Returning a value from a thread with std::packaged\_task<T>

```
template<typename Callable>
std::unique_future<std::result_of<Callable()>::type>
run_in_thread(Callable func)
{
    typedef std::result_of<Callable()>::type rtype;
    std::packaged_task<rtype> task(std::move(func));
    std::unique_future<rtype> res(task.get_future());
    std::thread(std::move(task)).detach();
    return std::move(res);
}
```

### Promises

- Value can come from any number of possible sources — e.g. first worker in pool to calculate result
- More explicit interface:
  - p.set\_value(some\_value)
  - p.set\_exception(some\_exception)

- Already some proposals for C++0x which have been retargeted to TR2
- shared\_mutex, upgrade\_mutex (from N2094)
- thread pools (from N2094, N2185, N2276)

- Provides a multiple-reader/single-writer mutex
- single writer:
  - m.lock()/m.unlock()
  - std::unique\_lock<shared\_mutex>
- multiple readers:
  - m.lock\_shared()/m.unlock\_shared()
  - shared\_lock<shared\_mutex>

#### upgrade\_mutex

- multiple readers + single upgrader / single writer
- The one and only upgrader can upgrade to a writer
  - Blocks until all readers have finished
  - Prevents other writers acquiring lock
  - Allows thread to rely on data read prior to upgrade
- Lock/unlock upgrader with:
  - m.lock\_upgrade()/m.unlock\_upgrade()
  - upgrade\_lock<upgrade\_mutex>
- Upgrade with:
  - m.unlock\_upgrade\_and\_lock()
  - Move-construction of an upgrade\_lock<upgrade\_mutex> to unique\_lock<upgrade\_mutex>
- Locks can be downgraded

# In boost 1.35.0, boost::shared\_mutex provides all this functionality.

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#### Thread Pools

- Universal agreement that we need to provide some kind of thread pool support.
- Exact API is not yet clear.
- N2094, N2185, N2276 provide distinct but similar APIs.
- Philipp Henkel has written a thread pool library that works with boost

- http://threadpool.sourceforge.net.

• Yigong Liu's Join library provides an alternative approach -- http://channel.sourceforge.net

# TR2 and beyond I

#### Thread Interruption

- Present in Boost 1.35.0
- Interrupt a thread by calling t.interrupt() on a thread object t.
- Thread throws thread\_interrupted exception at next *interruption point*
- Interruption points include condition\_variable::wait(), this\_thread::sleep() and interruption\_point()
- Interruption can be disabled with instances of disable\_interruption
- The thread\_interrupted exception can be caught: the thread can then be interrupted again

# TR2 and beyond II

- Thread-safe containers:
  - concurrent\_queue
  - concurrent\_stack
  - concurrent\_list
  - concurrent\_unordered\_map
- Parallel algorithms
  - parallel\_find
  - parallel\_sort
  - parallel\_accumulate
  - parallel\_for
- Intel TBB provides some of these

— http://threadingbuildingblocks.org/

# TR2 and beyond III

- Software Transactional Memory (STM) Allows for ACID guarantees in concurrent code, just like in databases
- OpenMP (http://www.openmp.org) A set of compiler directives to highlight code that should be run in parallel
- Auto-parallelisation in compilers
   A step beyond OpenMP compilers identify parallelizable regions automatically.

The current Intel compiler has basic support for this, with the -parallel command-line option.

### References and Further Reading

- The current C++0x working paper: N2588 http://www.open-std.org/jtc1/sc22/wg21/docs/ papers/2008/n2588.pdf
- The Boost 1.35.0 thread docs http://www.boost.org/doc/libs/1\_35\_0/doc/html/ thread.html
- The futures proposal: N2561 http://www.open-std.org/jtc1/sc22/wg21/docs/ papers/2008/n2561.html
- My book: C++ Concurrency in Action: Practical Multithreading, due to be published by Manning end of 2008/early 2009.